

Devicetree State of the Union



Agenda:

- Introduction
- System Devicetree in context, why we need it:
 - [Andrei Kholodnyi](#), Wind River, Principal Technologist, ROS2 TSC
- System Devicetree
 - [Stefano Stabellini](#), Xilinx, Principal Engineer
- Devicetree in Zephyr and other RTOS
 - [Kumar Gala](#), Linaro IOT and Embedded Group Technical Lead
- Devicetree in Standardized Linux & other OS boot
 - [Bill Mills](#), Linaro, Principal Technical Consultant
- Conclusion

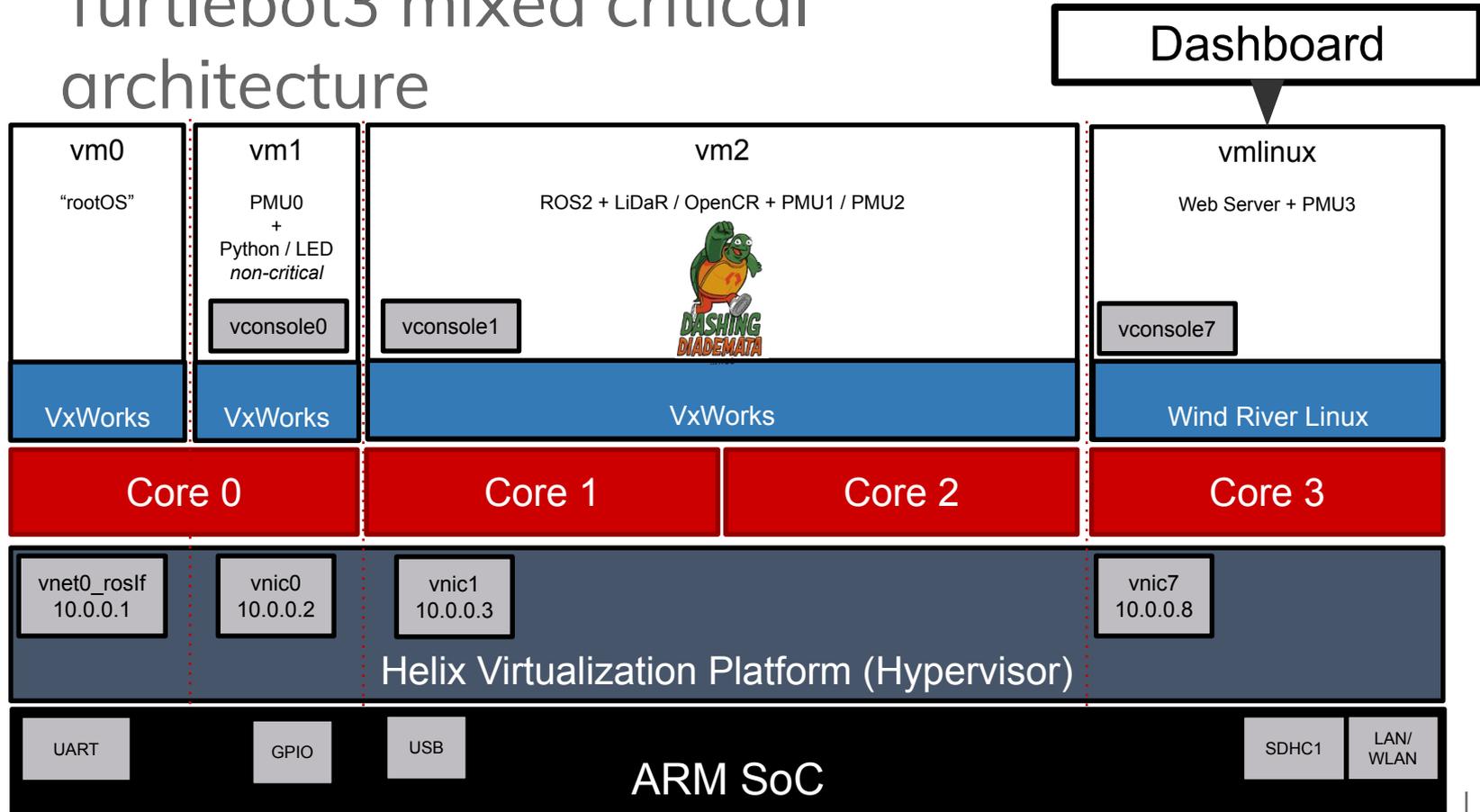
Introduction:

- What is Devicetree?
 - A structured data representation that describes SOC and board hardware
 - Started in the powerpc architecture, used for Arm in Linux since v3.1 in 2011
 - Supports many architectures:
 - arm, arm64, riscv, powerpc, mips, x86, microblaze, m68k, sh, arc, xtensa, nios, etc
 - Used in lots of projects:
 - Linux, U-Boot, TF-A, Xen, BSDs, Zephyr, VxWorks, Nucleus
- What is System Devicetree?
 - Describes the HW properties of the full system and how resources are allocated to different execution domains (e.g. Linux, XEN, RTOS)
- But ARM Server's use ACPI, why not just use that?
 - If ACPI meets your needs, use it
 - Devicetree is data only, ACPI is data + byte code. This kernel code remains in charge.
 - Devicetree is more flexible and nimble (blessing and a curse)
 - Devicetree can describe situations that ACPI cannot today
- Device Tree remains relevant today and in to the future
 - SystemReady IR is devicetree based to match current SOC practice
 - EBBR standardizes the boot flow for DT based systems

Need for System Devicetree



Turtlebot3 mixed critical architecture



Challenges to configure system with a hypervisor

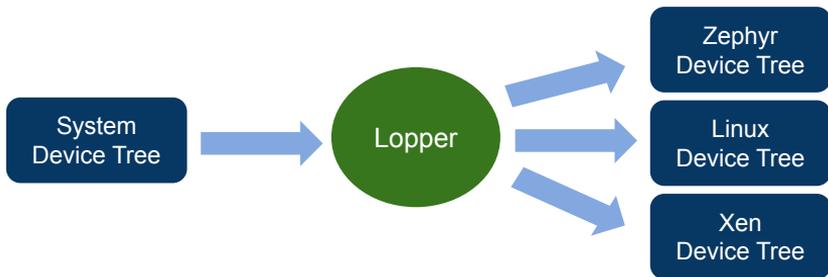
- System configuration is spreaded via different instances (hypervisor + RTOS + Linux)
 - CPUs
 - RAM
 - Network
 - Peripherals
- Configuration formats are not compatible
 - Proprietary Device Tree for the the hypervisor
 - VxWorks Device Tree (other RTOSes don't even have DT and values are hardcoded)
 - Linux Device Tree
- Difficult to check for consistency
- Modifications are needed in several places
- There are also AMP systems (e.g. with R5)

System Devicetree



System Device Tree

- Multi-View Multi-OS hardware description
 - *cpus,clusters*
 - *address-map*
 - *indirect-bus*
- AMP Configuration
 - execution domains (*openamp, domain*)
 - unique access to hardware resource (*access*)
- Lopper: prunes a S-DT down to traditional device trees



```
cpus_r5: cpus-cluster@0 {
    compatible = "cpus,cluster";

    address-map =
        <0xf1000000 &amba 0xf1000000 0xeb00000>,
        <0xf9000000 &amba_rpu 0xf9000000 0x10000>,
        <0x0 &memory 0x0 0x80000000>,
        <0x0 &tcm 0xFFE90000 0x10000>;

    cpu@0 {
    };

    cpu@1 {
    };
};

domains {
    domain@0 {
        compatible = "openamp, domain-v1";

        cpus = <&cpus_r5 0x3 0x80000001>;
        memory = <0x0 0x100000 0x0 0x400000>;

        access = <&can0 0x80000000>;
    };
};
```

OpenAMP RemoteProc

- Easier system-wide configuration in S-DT
 - dedicated resources (*access*)
 - shared resources (*include*)
- Lopper generates RemoteProc bindings from S-DT information
- Compatible with existing RemoteProc Bindings
 - No changes to existing software required

```
resource_group: resource_group@0 {
    compatible = "openamp,remoteproc-v1";
    memory = <0x0 0x3ed40000 0x0 0x4000
              0x0 0x3ed44000 0x0 0x4000
              0x0 0x3ed48000 0x0 0x100000
              0x0 0x3ed00000 0x0 0x40000>;
    access = <&tcm_0_a 0x0>;
};

openamp_a53 {
    compatible = "openamp,domain-v1";
    memory = <0x0 0x80000000 0x0 0x78000000
              0x8 0x0 0x0 0x80000000>;
    cpus = <&cpus_a53 0xf 0x2>;

    access = <&ipi_mailbox_rpu0 0x13>;
    include = <&resource_group 0x1>;
};

openamp_r5 {
    compatible = "openamp,domain-v1";
    memory = <0x0 0x0 0x0 0x20000000>;
    cpus = <&cpus_r5 0x2 0x80000000>;

    include = <&resource_group 0x0>;
};
```

Traditional RemoteProc Device Tree Output

```
r5fss@ff9a0000 {
    compatible = "xlnx,zynqmp-r5-remoteproc";
    #address-cells = <2>;
    #size-cells = <2>;
    ranges;
    reg = <0x0 0xff9a0000 0x0 0x10000>;
    xlnx,cluster-mode = <0>;

    r5f_0 {
        compatible = "xilinx,r5f";
        memory-region = <&elf_load0>,
            <&rpu0vdev0vring0>,
            <&rpu0vdev0vring1>,
            <&rpu0vdev0buffer>;
        sram = <&tcm_0a>, <&tcm_0b>;
        mboxes = <&ipi_mailbox_rpu0 0x0 &ipi_mailbox_rpu0 0x1>;
        mbox-names = "tx", "rx";
        power-domain = <0x7>;
    };
};
```

```
zynqmp_ipi1 {
    compatible = "xlnx,zynqmp-ipi-mailbox";
    interrupt-parent = <&gic>;
    interrupts = <0 33 4>;
    xlnx,ipi-id = <5>;
    #address-cells = <1>;
    #size-cells = <0>;

    ipi_mailbox_rpu0: mailbox@ff990600 {
        reg = <0xff990600 0x20>,
            <0xff990620 0x20>,
            <0xff9900c0 0x20>,
            <0xff9900e0 0x20>;
        reg-names = "local_request_region",
            "local_response_region",
            "remote_request_region",
            "remote_response_region";
        #mbox-cells = <1>;
        xlnx,ipi-id = <3>;
    };
};
```

System Device Tree Example

Full System Device Tree example [here](#), including:

- full hardware description of a Xilinx Versal board
- multiple domains

```
domains {

    #address-cells = <0x2>;
    #size-cells = <0x2>;

    linux: domain@0 {
        compatible = "openamp, domain-v1";
        cpus = <&cpus_a72 0x3 0x00000001>;
        memory = <0x0 0x500000 0x0 0x7fb00000>;

        access = <&mmc0 0x0 &ipi_mailbox_rpu0 0x13>;
        include = <&resource_group_0 0x1>;

        /* 1: block */
        firewallconfig-default = <1 0>;
    };

    zephyr: domain@1 {
        compatible = "openamp, domain-v1";
        cpus = <&cpus_r5 0x3 0x80000001>;
        memory = <0x0 0x100000 0x0 0x400000>;

        access = <&can0 0x0>;
        include = <&resource_group_0 0x0>;

        /* 1: block */
        firewallconfig-default = <1 0>;
    };
}
```

Next 6-12 months

- System Device Tree bindings specification
 - Bus Firewalls / System MMU bindings
 - Hypervisor bindings
 - YAML

Devicetree usage for RTOS/microcontroller



Devicetree for RTOS/Microcontroller

- Same reasons devicetree utilized by Linux, BSDs, Xen, u-boot, VxWorks, TF-A, ...
 - Vendor neutral hardware description
 - Convey hardware configuration of system to different software environments (RTOS, middleware, etc)
 - Ability to describe hardware configurations to different software contexts:
 - Asymmetric Multiprocessing
 - Trusted Execution Environments
 - Heterogeneous Systems (Linux + RTOS / MPU + MCU in same SoC)
- Unique concern for microcontroller software is footprint in memory usage, code and data size
 - MCUs having 256K of flash, 64k of memory (BBC Microbit)
 - Devicetree blob - 6k
 - Libfdt - 48k
- Utilize devicetree source as a common data representation and generate code specific to software environment need

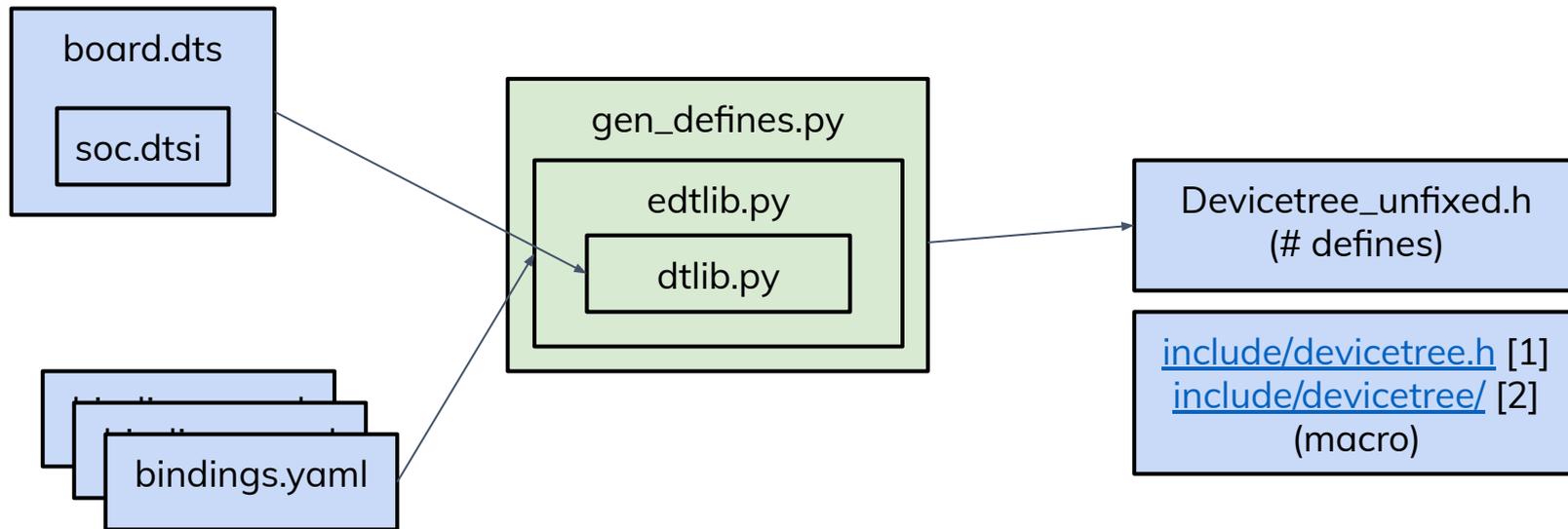
Devicetree codegen RTOS/Microcontroller

- Zephyr RTOS was utilizing Kconfig to describe hardware instance info
- Linaro introduced devicetree to Zephyr in version 1.7 / Jan 2017
- Tooling based around python - need portability across Linux/MacOS X/Windows
- Several iterations of tooling and what code generated looks like
- As of Zephyr 2.5 (Feb 12, 2021) - 3 python scripts for code generation:
 - scripts try not to be Zephyr specific:
 - [dtlib.py](#) - pure python devicetree source parser
 - [edtlib.py](#) - utilizes dtlib.py + devicetree bindings to provide device data model view (concepts like registers, interrupts, gpios, clocks, etc.)
 - [gen_defines.py](#) - utilizes edtlib.py to generate defines
- Generate a set of #defines + macros that are able to extract devicetree data

```
uart0: uart@4006a000 {
    compatible = "nxp,kinetis-uart";
    reg = < 0x4006a000 0x1000 >;
};
```

- **Example** - `DT_REG_ADDR(DT_NODELABEL(uart0)) /* 0x4006a000 */`

Zephyr code generation workflow



[1] <https://github.com/zephyrproject-rtos/zephyr/blob/master/include/devicetree.h>

[2] <https://github.com/zephyrproject-rtos/zephyr/tree/master/include/devicetree>

What's next for RTOS/MCU and devicetree

- Expanding into uses cases that need system devicetree:
 - TFM (Trusted Firmware-M) and ability to utilize system devicetree to describe secure and non-secure resource partitioning and have TFM and Zephyr utilize output
- Work with system devicetree group on ensuring bindings defined for defining contexts and partition handle MCU hardware and usecases
 - Various hardware protection units for MCUs, partitioning of device resources, memory, flash, OpenAMP/remoteproc/rpmsg, etc.
- Work with looper to utilize a common python scripts/libraries to build tools around:
 - Define a common data view of devicetree
 - Have a common python library that can read and write devicetree(s)
- Example system devicetree workflow for Trusted Firmware-M

Devicetree in Standard Linux Boot



Devicetree in Linux and other OS Boot

- EBBR and Arm SystemReady IR
 - Standardizing boot in traditional embedded Linux form factors
 - Allow Separation of OS/Hypervisors layer from Firmware layer
 - Devicetree is an accepted part of this profile
- Cortex-A profile Operating systems and Hypervisors
 - Linux, BSDs, Xen, etc
- The existing boot flow patterns are too varied and too platform specific
 - No clear distinction between Firmware layer and OS layer
 - Varies from SOC vendor to SOC vendor
 - Varies year to year for Firmware / OS pairs from the same SOC vendor
- DT must co-exist with Secure and Measured boot
 - If doing secure boot, all DTB data must be verified by signature
 - If doing measured boot, all DTB data must be measured

Who owns the DTB data?

- Everyone agrees the answer is obvious
 - They just disagree on which answer that is
- Two fundamental models for ownership of DTB used by OS
 - DTB is supplied by Firmware
 - Simple and better in theory, more standardized
 - New HW should just work, no need for OS to know exact model
 - Harder to test and enforce
 - Not current practice
 - DTB is supplied by OS
 - Exact match between DTB and kernel version
 - Matches current common practice
 - OS works only with boards it knows
 - Current Plan: Hybrid
 - DTB is supplied by firmware but OS boot manager can override if it knows better
 - OS works great with boards it knows and OK or great for others based on ABI stability
 - Still need ABI and testing

Device Tree Evolution: Current Activity

- Support for DT overlay source in the upstream kernel
 - Boot loader applied overlays are an existing pattern
 - **However**, up till now, no place to call upstream
 - Fixing this is 5.12 (started) and rest in 5.13 (or later 5.12-rcN)
 - Build and test apply during kernel build
 - Future work to do more verification on overlay source
- UEFI DT Fixup API for use by OS Boot managers (Grub etc)
 - The DT from OS has a problem when using an OS Boot manager
 - Grub can load a new DTB but does not know how to fixup MAC addr's, SN, etc
 - New UEFI API defined so Grub can call back to firmware to do fixups
 - Implemented in U-boot master
- Proof of Concept build for DT from firmware
 - Firmware: TFA, U-boot, OP-TEE, DTB for OS
 - UEFI based boot w/ or w/o Grub
 - QEMU based but built in a way that does not rely on QEMU DTB magic
 - Example to be used for real HW

Device Tree Evolution: Next 6 to 12 months

- Enable DT from OS model in a standard compatible fashion
 - Complete the Grub based “DTB load w/ firmware fixup” boot model
- Begin testing DT ABI stability
 - Begin testing kernel tip against DTB from previous kernels
 - Start with select boards
 - Start with DTBs from previous 1 or 2 LTS’s
 - Later expand to more boards and maybe new DTB with old kernel
- Reporting tool to check cross project DT synchronization
 - One DT repo “to rule them all” has been rejected
 - Instead cross check project DT source synchronization
 - Tool for project maintainers
- Define new DTB file/memory format
 - Many suggestions for new FDT formats have accumulated over the past few years
 - Some or all of the below will make it:
 - Metadata: Source files hash, build date, **signature of data**
 - Better overlay support: delete node & property, format clean-up
 - Direct support of segmented FDT (base + overlays + fixup)
 - Size reduction: string deduplication or elimination, compression

Conclusion



Conclusion:

- Devicetree is an important technology that is driving innovation in many areas
 - Heterogeneous Systems
 - RTOS and High Level OS (Linux etc)
 - Hypervisors
 - Secure OS and Firmware
- Devicetree is driving more standardization
 - Devicetree is more nimble and wider scope than other technologies
 - **But** we need stability to drive many use cases like OS boot (Linux, etc)

No Q&A in this session

Please stay for the Birds of a Feather session directly following this one

All speakers in this session and others will be available for questions

Thank you

Accelerating deployment in the Arm Ecosystem

